

Review.



Theory: If you drink alcohol you must be at least 18.

Which cards do you turn over?

Drink Alcohol \implies " ≥ 18 "

" < 18 " \implies Don't Drink Alcohol. Contrapositive.

(A) (B) (C) and/or (D)?

Propositional Forms: $\wedge, \vee, \neg, P \implies Q \equiv \neg P \vee Q$.

Truth Table. Putting together identities. (E.g., cases, substitution.)

Predicates, $P(x)$, and quantifiers. $\forall x, P(x)$.

DeMorgan's: $\neg(P \vee Q) \equiv \neg P \wedge \neg Q$. $\neg \forall x, P(x) \equiv \exists x, \neg P(x)$.

Quick Background, Notation and Definitions!

Integers closed under addition.

$$a, b \in \mathbb{Z} \implies a + b \in \mathbb{Z}$$

$a|b$ means "a divides b".

$2|4$? Yes! Since for $q = 2$, $4 = (2)2$.

$7|23$? No! No q where true.

$4|2$? No!

$2|-4$? Yes! Since for $q = 2$, $-4 = (-2)2$.

Formally: for $a, b \in \mathbb{Z}$, $a|b \iff \exists q \in \mathbb{Z}$ where $b = aq$.

$3|15$ since for $q = 5$, $15 = 3(5)$.

A natural number $p > 1$, is **prime** if it is divisible only by 1 and itself.

A number x is even if and only if $2|x$, or $x = 2k$ for $x, k \in \mathbb{Z}$.

A number x is odd if and only if $x = 2k + 1$

CS70: Lecture 2. Outline.

Today: Proofs!!!

1. By Example.
2. Direct. (Prove $P \implies Q$.)
3. by Contraposition (Prove $P \implies Q$)
4. by Contradiction (Prove P .)
5. by Cases

If time: discuss induction.

Divides.

$a|b$ means

- (A) There exists $k \in \mathbb{Z}$, with $a = kb$.
- (B) There exists $k \in \mathbb{Z}$, with $b = ka$.
- (C) There exists $k \in \mathbb{N}$, with $b = ka$.
- (D) There exists $k \in \mathbb{Z}$, with $k = ab$.
- (E) a divides b

Incorrect: (C) sufficient not necessary. (A) Wrong way. (D) the product is an integer.

Correct: (B) and (E).

Last time: Existential statement.

How to prove existential statement?

Give an example. (Sometimes called "proof by example.")

Theorem: $(\exists x \in \mathbb{N})(x = x^2)$

Pf: $0 = 0^2 = 0$

□

Often used to disprove claim.

Homework.

Direct Proof.

Theorem: For any $a, b, c \in \mathbb{Z}$, if $a|b$ and $a|c$ then $a|(b - c)$.

Proof: Assume $a|b$ and $a|c$

$$b = aq \text{ and } c = aq' \text{ where } q, q' \in \mathbb{Z}$$

$$b - c = aq - aq' = a(q - q') \text{ Done?}$$

$(b - c) = a(q - q')$ and $(q - q')$ is an integer so by definition of divides $a|(b - c)$ □

Works for $\forall a, b, c$?

Argument applies to every $a, b, c \in \mathbb{Z}$.

Used distributive property and definition of divides.

Direct Proof Form:

Goal: $P \implies Q$

Assume P .

...

Therefore Q .

Another direct proof.

Let D_3 be the 3 digit natural numbers.

Theorem: For $n \in D_3$, if the alternating sum of digits of n is divisible by 11, then $11|n$.

$$\forall n \in D_3, (11 | \text{alt. sum of digits of } n) \implies 11 | n$$

Examples:

$n = 121$ Alt Sum: $1 - 2 + 1 = 0$. Divis. by 11. As is 121.

$n = 605$ Alt Sum: $6 - 0 + 5 = 11$ Divis. by 11. As is $605 = 11(55)$

Proof: For $n \in D_3$, $n = 100a + 10b + c$, for some a, b, c .

Assume: Alt. sum: $a - b + c = 11k$ for some integer k .

Add $99a + 11b$ to both sides.

$$100a + 10b + c = 11k + 99a + 11b = 11(k + 9a + b)$$

Left hand side is n , $k + 9a + b$ is integer. $\implies 11 | n$. □

Direct proof of $P \implies Q$:

Assumed P : $11 | a - b + c$. Proved Q : $11 | n$.

Proof by Contraposition

Thm: For $n \in \mathbb{Z}^+$ and $d | n$. If n is odd then d is odd.

$n = kd$ and $n = 2k' + 1$ for integers k, k' .

what do we know about d ?

Goal: Prove $P \implies Q$.

Assume $\neg Q$

...and prove $\neg P$.

Conclusion: $\neg Q \implies \neg P$ equivalent to $P \implies Q$.

Proof: Assume $\neg Q$: d is even. $d = 2k$.

$d | n$ so we have

$$n = qd = q(2k) = 2(kq)$$

n is even. $\neg P$ □

The Converse

Thm: $\forall n \in D_3, (11 | \text{alt. sum of digits of } n) \implies 11 | n$

Is converse a theorem?

$\forall n \in D_3, (11 | n) \implies (11 | \text{alt. sum of digits of } n)$

Yes? No?

Another Contraposition...

Lemma: For every n in \mathbb{N} , n^2 is even $\implies n$ is even. ($P \implies Q$)

n^2 is even, $n^2 = 2k$, ... $\sqrt{2k}$ even?

Proof by contraposition: ($P \implies Q$) \equiv ($\neg Q \implies \neg P$)

P = ' n^2 is even' $\neg P$ = ' n^2 is odd'

Q = ' n is even' $\neg Q$ = ' n is odd'

Prove $\neg Q \implies \neg P$: n is odd $\implies n^2$ is odd.

$$n = 2k + 1$$

$$n^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1.$$

$$n^2 = 2l + 1 \text{ where } l \text{ is a natural number..}$$

... and n^2 is odd!

$\neg Q \implies \neg P$ so $P \implies Q$ and ... □

Another Direct Proof.

Theorem: $\forall n \in D_3, (11 | n) \implies (11 | \text{alt. sum of digits of } n)$

Proof: Assume $11 | n$.

$$n = 100a + 10b + c = 11k \implies$$

$$99a + 11b + (a - b + c) = 11k \implies$$

$$a - b + c = 11k - 99a - 11b \implies$$

$$a - b + c = 11(k - 9a - b) \implies$$

$$a - b + c = 11\ell \text{ where } \ell = (k - 9a - b) \in \mathbb{Z}$$

That is $11 | \text{alternating sum of digits}$. □

Note: similar proof to other direction. In this case every \implies is \iff

Often works with arithmetic properties ...

...not when multiplying by 0.

We have.

Theorem: $\forall n \in \mathbb{N}', (11 | \text{alt. sum of digits of } n) \iff (11 | n)$

Proof by Obfuscation.



ob·fus·ca·tion

/,äbfə'skāSH(ə)n/

noun

noun: **obfuscation**; plural noun: **obfuscations**

the action of making something obscure, unclear, or unintelligible.
"when confronted with sharp questions they resort to obfuscation"

Proof by contradiction: form

Theorem: $\sqrt{2}$ is irrational.

Must show: For every $a, b \in \mathbb{Z}$, $(\frac{a}{b})^2 \neq 2$.

A simple property (equality) should always “not” hold.

Proof by contradiction:

Theorem: P .

$\neg P \implies P_1 \dots \implies R$

$\neg P \implies Q_1 \dots \implies \neg R$

$\neg P \implies R \wedge \neg R \equiv \text{False}$

or $\neg P \implies \text{False}$

Contrapositive of $\neg P \implies \text{False}$ is $\text{True} \implies P$.

Theorem P is true. And proven. \square

Contradiction

Theorem: $\sqrt{2}$ is irrational.

Assume $\neg P$: $\sqrt{2} = a/b$ for $a, b \in \mathbb{Z}$.

Reduced form: a and b have no common factors.

$$\sqrt{2}b = a$$

$$2b^2 = a^2 = 4k^2$$

a^2 is even $\implies a$ is even.

$a = 2k$ for some integer k

$$b^2 = 2k^2$$

b^2 is even $\implies b$ is even.

a and b have a common factor. Contradiction. \square

Proof by contradiction: example

Theorem: There are infinitely many primes.

Proof:

► Assume finitely many primes: p_1, \dots, p_k .

► Consider number

$$q = (p_1 \times p_2 \times \dots \times p_k) + 1.$$

► q cannot be one of the primes as it is larger than any p_i .

► q has prime divisor p (“ $p > 1$ ” = R) which is one of p_i .

► p divides both $x = p_1 \cdot p_2 \cdot \dots \cdot p_k$ and q , and divides $q - x$,

► $\implies p \mid (q - x) \implies p \leq (q - x) = 1$.

► so $p \leq 1$. (**Contradicts R .**)

The original assumption that “the theorem is false” is false, thus the theorem is proven. \square

Product of first k primes..

Did we prove?

- “The product of the first k primes plus 1 is prime.”
- No.
- The chain of reasoning started with a false statement.

Consider example..

- $2 \times 3 \times 5 \times 7 \times 11 \times 13 + 1 = 30031 = 59 \times 509$
- There is a prime *in between* 13 and $q = 30031$ that divides q .
- Proof assumed no primes *in between* p_k and q .
As it assumed the only primes were the first k primes.

Poll: Odds and evens.

x is even, y is odd.

Even numbers are divisible by 2.

Which are even?

(A) x^3 **Even: $(2k)^3 = 2(4k^3)$**

(B) y^3

(C) $x + 5x$ **Even: $2k + 5(2k) = 2(k + 5k)$**

(D) xy **Even: $2(ky)$.**

(E) xy^5 **Even: $2(ky^5)$.**

(F) $x + y$

A, C, D, E all contain a factor of 2.

E.g., $x = 2k$, $x^3 = 8k = 2(4k)$ and is even.

y^3 . Odd?

$$y = (2k + 1). \quad y^3 = 8k^3 + 24k^2 + 24k + 1 = 2(4k^3 + 12k^2 + 12k) + 1.$$

Odd times an odd? Odd.

Any power of an odd number? Odd.

Idea: $(2k + 1)^n$ has terms

(a) with the last term being 1

(b) and all other terms having a multiple of $2k$.

Proof by cases.

Theorem: $x^5 - x + 1 = 0$ has no solution in the rationals.

Proof: First a lemma..

Lemma: If x is a solution to $x^5 - x + 1 = 0$ and $x = a/b$ for $a, b \in \mathbb{Z}$, then both a and b are even.

Reduced form $\frac{a}{b}$: a and b can't both be even! + Lemma \implies no rational solution. \square

Proof of lemma: Assume a solution of the form a/b .

$$\left(\frac{a}{b}\right)^5 - \frac{a}{b} + 1 = 0$$

Multiply by b^5 ,

$$a^5 - ab^4 + b^5 = 0$$

Case 1: a odd, b odd: odd - odd + odd = even. **Not possible.**

Case 2: a even, b odd: even - even + odd = even. **Not possible.**

Case 3: a odd, b even: odd - even + even = even. **Not possible.**

Case 4: a even, b even: even - even + even = even. **Possible.**

The fourth case is the only one possible, so the lemma follows. \square

Proof by cases.

Theorem: There exist irrational x and y such that x^y is rational.

Let $x = y = \sqrt{2}$.

Case 1: $x^y = \sqrt{2}^{\sqrt{2}}$ is rational. Done!

Case 2: $\sqrt{2}^{\sqrt{2}}$ is irrational.

► New values: $x = \sqrt{2}^{\sqrt{2}}$, $y = \sqrt{2}$.

►

$$x^y = \left(\sqrt{2}^{\sqrt{2}} \right)^{\sqrt{2}} = \sqrt{2}^{\sqrt{2} \cdot \sqrt{2}} = \sqrt{2}^2 = 2.$$

Thus, we have irrational x and y with a rational x^y (i.e., 2).

One of the cases is true so theorem holds. □

Question: Which case holds? Don't know!!!

Poll: proof review.

Which of the following are (certainly) true?

(A) $\sqrt{2}$ is irrational.

(B) $\sqrt{2}^{\sqrt{2}}$ is rational.

(C) $\sqrt{2}^{\sqrt{2}}$ is rational or it isn't.

(D) $(2^{\sqrt{2}})^{\sqrt{2}}$ is rational.

(A),(C),(D)

(B) I don't know.

Be careful.

Theorem: $3 = 4$

Proof: Assume $3 = 4$.

Start with $12 = 12$.

Divide one side by 3 and the other by 4 to get
 $4 = 3$.

By commutativity theorem holds. □

What's wrong?

Don't assume what you want to prove!

Be really careful!

Theorem: $1 = 2$

Proof: For $x = y$, we have

$$(x^2 - xy) = x^2 - y^2$$

$$x(x - y) = (x + y)(x - y)$$

$$x = (x + y)$$

$$x = 2x$$

$$1 = 2$$

Poll: What is the problem? □

(A) Assumed what you were proving.

(B) No problem. It's fine.

(C) $x - y$ is zero.

(D) Can't multiply by zero in a proof.

Dividing by zero is no good. Multiplying by zero is weirdly cool!

Also: Multiplying inequalities by a negative.

$P \implies Q$ does not mean $Q \implies P$.

Summary: Note 2.

Direct Proof:

To Prove: $P \implies Q$. Assume P . Prove Q .

$a|b$ and $a|c \implies a|(b+c)$.

By Contraposition:

To Prove: $P \implies Q$ Assume $\neg Q$. Prove $\neg P$.

n^2 is odd $\implies n$ is odd. $\equiv n$ is even $\implies n^2$ is even.

By Contradiction:

To Prove: P Assume $\neg P$. Prove False.

$\sqrt{2}$ is rational.

$\sqrt{2} = \frac{a}{b}$ with no common factors....

By Cases: informal.

Universal: show that statement holds in all cases.

Existence: used cases where one is true.

Either $\sqrt{2}$ and $\sqrt{2}$ worked.

or $\sqrt{2}$ and $\sqrt{2}^{\sqrt{2}}$ worked.

Careful when proving!

Don't assume the theorem. Divide by zero. Watch converse. ...

CS70: Note 3. Induction!

Poll. What's the biggest number?

(A) 100

(B) 101

(C) $n+1$

(D) infinity.

(E) This is about the "recursive leap of faith."